

## Operating Systems and Networks

### Networks Part 2: Physical Layer

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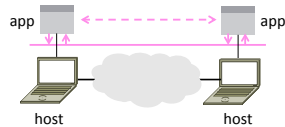
## Overview

- Important concepts from last lecture
  - Statistical multiplexing, statistical multiplexing gain
  - OSI 7 layer model, interfaces, protocols
  - Encapsulation, demultiplexing
- This lecture
  - Socket programming overview
  - Physical layer
- Online lecture videos: <http://computernetworks5e.org>

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## Network-Application Interface

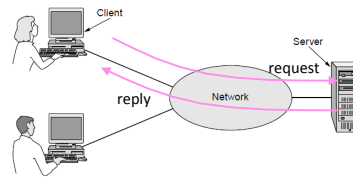
- Defines how apps use the network
  - Lets apps talk to each other via hosts; hides the details of the network



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## Motivating Application

- Simple client-server connection setup



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## Motivating Application (2)

- Simple client-server connection setup
  - Client app sends a request to server app
  - Server app returns a (longer) reply
- This is the basis for many apps!
  - File transfer: send name, get file (§6.1.4)
  - Web browsing: send URL, get page
  - Echo: send message, get it back
- Let's see how to write this app ...

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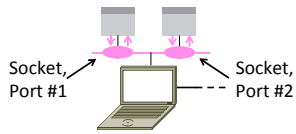
## Socket API

- Simple abstraction to use the network
  - The network service API used to write all Internet applications
  - Part of all major OSes and languages; originally Berkeley (Unix) ~1983
- Supports two kinds of network services
  - Streams: reliably send a stream of bytes
  - Datagrams: unreliably send separate messages. (Ignore for now.)

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## Socket API (2)

- Sockets let applications attach to the local network at different ports



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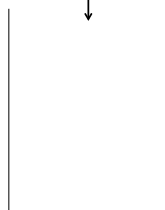
## Socket API (3)

Primitive	Meaning
SOCKET	Create a new communication endpoint
BIND	Associate a local address with a socket
LISTEN	Announce willingness to accept connections; give queue size
ACCEPT	Passively wait for an incoming connection
CONNECT	Actively attempt to establish a connection
SEND	Send some data over the connection
RECEIVE	Receive some data from the connection
CLOSE	Release the connection

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## Using Sockets

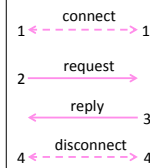
Client (host 1)    Time    Server (host 2)



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## Using Sockets (2)

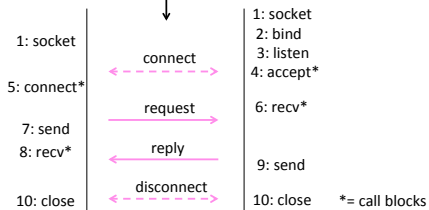
Client (host 1)    Time    Server (host 2)



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## Using Sockets (3)

Client (host 1)    Time    Server (host 2)



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## Client Program (outline)

```

socket()           // make socket
getaddrinfo()     // server and port name
                  // www.example.com:80
connect()         // connect to server [block]
...
send()            // send request
recv()            // await reply [block]
...
                  // do something with data!
close()           // done, disconnect
    
```

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## Server Program (outline)

```

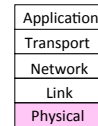
socket() // make socket
getaddrinfo() // for port on this host
bind() // associate port with socket
listen() // prepare to accept connections
accept() // wait for a connection [block]
...
recv() // wait for request
...
send() // send the reply
close() // eventually disconnect

```

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## Where we are in the Course

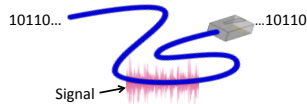
- Beginning to work our way up starting with the Physical layer



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## Scope of the Physical Layer

- Concerns how signals are used to transfer message bits over a link
  - Wires etc. carry analog signals
  - We want to send digital bits



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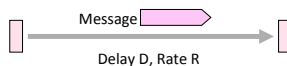
## Topics

- Properties of media
  - Wires, fiber optics, wireless
- Simple signal propagation
  - Bandwidth, attenuation, noise
- Modulation schemes
  - Representing bits, noise
- Fundamental limits
  - Nyquist, Shannon

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## Simple Link Model

- We'll end with an abstraction of a physical channel
  - Rate (or bandwidth, capacity, speed) in bits/second
  - Delay or Latency in seconds, related to length



- Other important properties:
  - Whether the channel is broadcast, and its error rate

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## Message Latency

- Latency L: delay to send a message over a link
  - Transmission delay: time to put M-bit message "on the wire"
 
$$T\text{-delay} = M \text{ (bits)} / \text{Rate (bits/sec)} = M/R \text{ seconds}$$
  - Propagation delay: time for bits to propagate across the wire
 
$$P\text{-delay} = \text{Length} / \text{speed of signals} = \text{Length} / \frac{1}{3}c = D \text{ seconds}$$
  - Combining the two terms we have:  $L = M/R + D$

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## Metric Units

- The main prefixes we use:

Prefix	Exp.	prefix	exp.
K(ilo)	$10^3$	m(illi)	$10^{-3}$
M(ega)	$10^6$	$\mu$ (micro)	$10^{-6}$
G(iga)	$10^9$	n(ano)	$10^{-9}$

- Use powers of 10 for rates, 2 for storage or data size
  - 1 Mbps = 1,000,000 bps, 1 KB =  $2^{10}$  bytes
- "B" is for bytes, "b" is for bits

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## Latency Examples

- "Dialup" with a telephone modem:
  - D = 5 ms, R = 56 kbps, M = 1250 bytes
- Broadband cross-country link:
  - D = 50 ms, R = 10 Mbps, M = 1250 bytes


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## Latency Examples (2)

- "Dialup" with a telephone modem:
  - D = 5 ms, R = 56 kbps, M = 1250 bytes
  - $L = 5 \text{ ms} + (1250 \times 8) / (56 \times 10^3) \text{ sec} = 184 \text{ ms!}$
- Broadband cross-country link:
  - D = 50 ms, R = 10 Mbps, M = 1250 bytes
  - $L = 50 \text{ ms} + (1250 \times 8) / (10 \times 10^6) \text{ sec} = 51 \text{ ms}$
- A long link or a slow rate means high latency
  - Often, one delay component dominates

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## Bandwidth-Delay Product

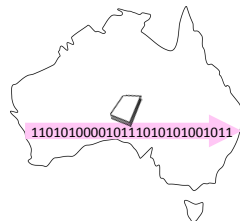
- Messages take space on the wire!
- 
- The amount of data in flight is the bandwidth-delay (BD) product

$$BD = R \times D$$
    - Measure in bits, or in messages
    - Small for LANs, big for "long fat" pipes

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## Bandwidth-Delay Example

- Fiber at home, cross-country
  - R=40 Mbps, D=50 ms
  - $BD = 40 \times 10^6 \times 50 \times 10^{-3} \text{ bits}$
  - = 2000 Kbit
  - = 250 KB
- That's quite a lot of data "in the network"!



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## How "Long" is a Bit?

- Interesting trivia: how "long" is the representation of a bit on a wire?
- Considering a fiber optic cable
  - Signal propagation speed: 200'000'000 m/s
  - Sending rate: 1Gbps  $\rightarrow$  duration of sending one bit: 1ns
  - Bit "length":  $1 \text{ ns} \times 200'000'000 \text{ m/s} = 0.2 \text{ m}$
  - "Length" of a 1Kb packet:  $0.2 \text{ m} \times 8 \times 2^{10} = 1.6 \text{ km}$

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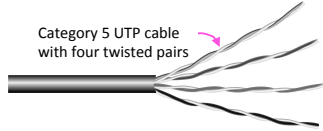
### Types of Media (§2.2, 2.3)

- Media propagate signals that carry bits of information
- We'll look at some common types:
  - Wires
  - Fiber (fiber optic cables)
  - Wireless

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### Wires – Twisted Pair

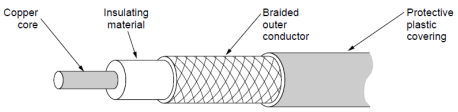
- Very common; used in LANs and telephone lines
  - Twists can reduce radiated signal or reduce effect of external interference signal



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### Wires – Coaxial Cable

- Also common. Better shielding for better performance

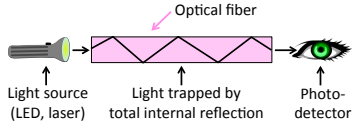


- Other kinds of wires too: e.g., electrical power (§2.2.4)

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### Fiber

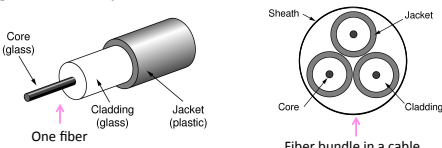
- Long, thin, pure strands of glass
  - Enormous bandwidth (high speed) over long distances



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### Fiber (2)


- Two varieties: multi-mode (shorter links, cheaper) and single-mode (up to ~100 km)



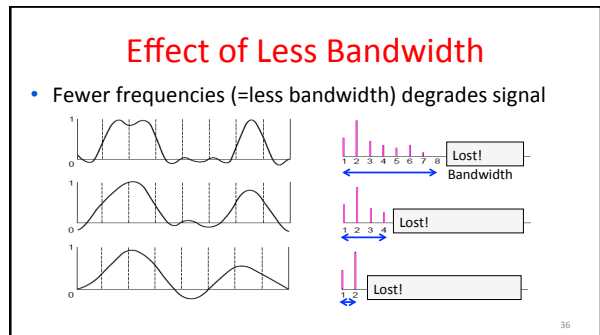
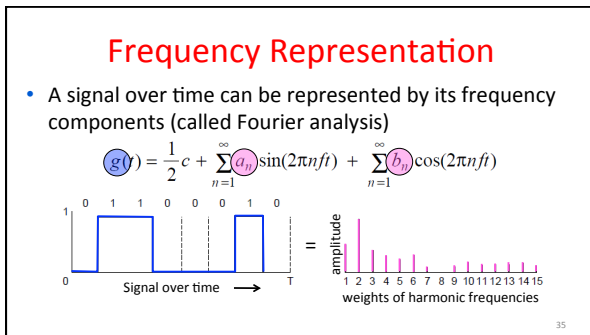
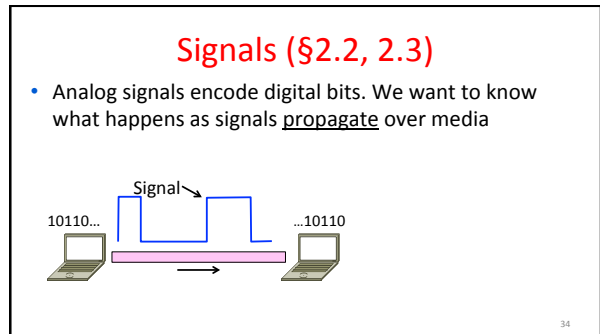
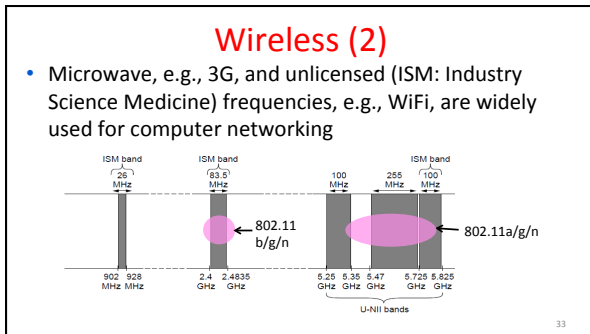
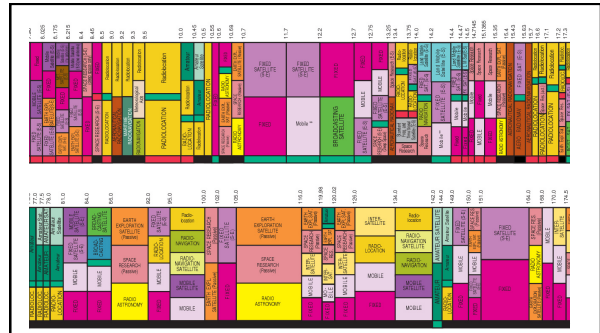
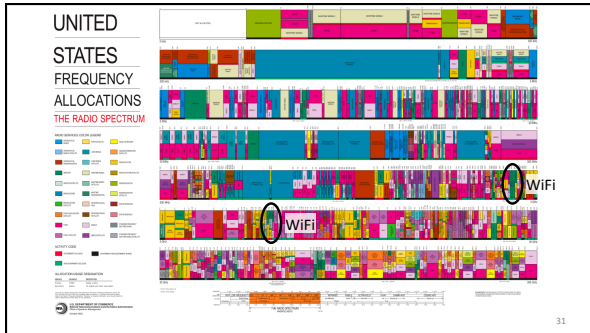
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### Wireless

- Sender radiates signal over a region
  - In many directions, unlike a wire, to potentially many receivers
  - Nearby signals (same freq.) interfere at a receiver; need to coordinate use



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## Signals over a Wire

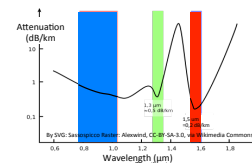
- What happens to a signal as it passes over a wire?
  1. The signal is delayed (propagates at  $\frac{2}{3}c$ )
  2. The signal is attenuated
  3. Frequencies above a cutoff are highly attenuated
  4. Noise is added to the signal (later, causes errors)

EE: Bandwidth = width of frequency band, measured in Hz  
 CS: Bandwidth = information carrying capacity, in bits/sec

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## Signals over Fiber

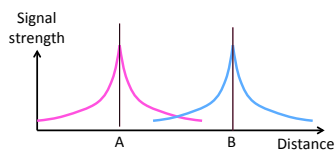
- Light propagates with very low loss in three very wide frequency bands
  - Use a carrier to send information



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## Signals over Wireless

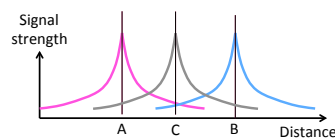
- Travel at speed of light, spread out and attenuate faster than  $1/\text{dist}^2$



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## Signals over Wireless (2)

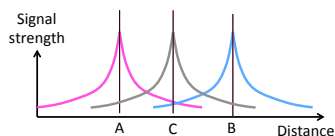
- Multiple signals on the same frequency interfere at a receiver



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## Signals over Wireless (3)

- Interference leads to notion of spatial reuse (of same freq.)



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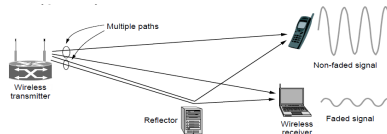
## Signals over Wireless (4)

- Various other effects too!
  - Wireless propagation is complex, depends on environment
- Some key effects are highly frequency dependent
  - E.g., multipath at microwave frequencies

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## Wireless Multipath

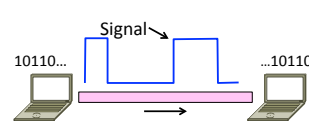
- Signals bounce off objects and take multiple paths
  - Some frequencies attenuated at receiver, varies with location
  - Messes up signal; handled with sophisticated methods



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## Modulation (§2.5)

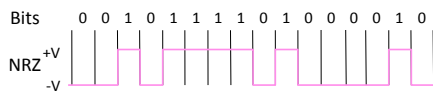
- We've talked about signals representing bits. How, exactly?
  - This is the topic of modulation



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## A Simple Modulation

- Let a high voltage (+V) represent a 1, and low voltage (-V) represent a 0
  - This is called NRZ (Non-Return to Zero)



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## Many Other Schemes

- Can use more signal levels, e.g., 4 levels is 2 bits per symbol
- Practical schemes are driven by engineering considerations
  - E.g., clock recovery

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## Clock Recovery

- Um, how many zeros was that?
  - Receiver needs frequent signal transitions to decode bits

1 0 0 0 0 0 0 0 0 ... 0

- Several possible designs
  - E.g., Manchester coding and scrambling (§2.5.1)

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## Clock Recovery – 4B/5B

- Map every 4 data bits into 5 code bits without long runs of zeros
  - 0000 → 11110, 0001 → 01001, 1110 → 11100, ... 1111 → 11101
  - Has at most 3 zeros in a row
  - Also invert signal level on a 1 to break up long runs of 1s (called NRZI, §2.5.1)

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### Clock Recovery – 4B/5B (2)

- 4B/5B code for reference:
  - 0000 → 11110, 0001 → 01001, 1110 → 11100, ... 1111 → 11101

• Message bits: 1 1 1 1 0 0 0 0 0 0 1

Coded Bits: 1 1 1 0 1 1 1 1 0 0 1 0 0 1



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### Passband Modulation

- What we have seen so far is baseband modulation for wires
  - Signal is sent directly on a wire
- These signals do not propagate well on fiber / wireless
  - Need to send at higher frequencies
- Passband modulation carries a signal by modulating a carrier

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### Passband Modulation (2)

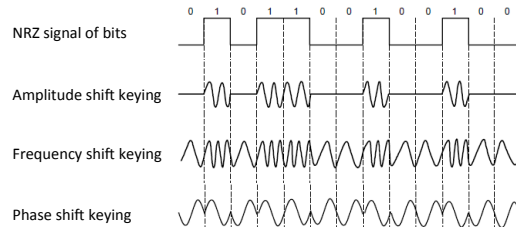
- Carrier is simply a signal oscillating at a desired frequency:



- We can modulate it by changing:
  - Amplitude, frequency, or phase

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### Passband Modulation (3)



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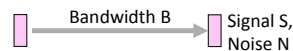
### Fundamental Limits (§2.1)

- How rapidly can we send information over a link?
  - Nyquist limit (~1924)
  - Shannon capacity (1948)
- Practical systems are devised to approach these limits

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### Key Channel Properties

- The bandwidth (B), signal strength (S), and noise strength (N)
  - B limits the rate of transitions
  - S and N limit how many signal levels we can distinguish

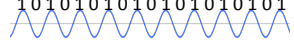


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### Nyquist Limit

- The maximum symbol rate is  $2B$

101010101010101010101




- Thus if there are  $V$  signal levels, ignoring noise, the maximum bit rate is:

$$R = 2B \log_2 V \text{ bits/sec}$$

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### Claude Shannon (1916-2001)

- Father of information theory
  - "A Mathematical Theory of Communication", 1948
- Fundamental contributions to digital computers, security, and communications

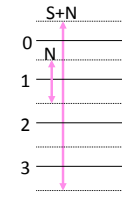


Electromechanical mouse that "solves" mazes!  
Credit: Courtesy MIT Museum

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### Shannon Capacity

- How many levels we can distinguish depends on  $S/N$ 
  - Or SNR, the Signal-to-Noise Ratio
  - Note noise is random, hence some errors
- SNR given on a log-scale in deciBels:
  - $SNR_{dB} = 10 \log_{10}(S/N)$



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### Shannon Capacity (2)

- Shannon limit is for capacity ( $C$ ), the maximum information carrying rate of the channel:

$$C = B \log_2(1 + S/N) \text{ bits/sec}$$

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### Wired/Wireless Perspective

- Wires and Fiber** Engineer SNR for data rate
  - Engineer link to have requisite SNR and  $B$
  - Can fix data rate
- Wireless** Adapt data rate to SNR
  - Given  $B$ , but SNR varies greatly, e.g., up to 60 dB!
  - Can't design for worst case, must adapt data rate

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### Putting it all together – DSL

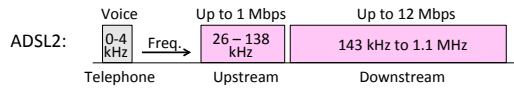
- DSL (Digital Subscriber Line, see §2.6.3) is widely used for broadband; many variants offer 10s of Mbps
  - Reuses twisted pair telephone line to the home; it has up to ~2 MHz of bandwidth but uses only the lowest ~4 kHz

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## DSL (2)

- DSL uses passband modulation (called OFDM §2.5.1)
  - Separate bands for upstream and downstream (larger)
  - Modulation varies both amplitude and phase (called QAM)
  - High SNR, up to 15 bits/symbol, low SNR only 1 bit/symbol



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