## **Operating Systems and Networks**

## Networks Part 2: Physical Layer

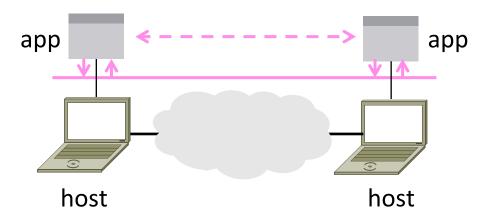
Adrian Perrig Network Security Group ETH Zürich

#### Overview

- Important concepts from last lecture
  - Statistical multiplexing, statistical multiplexing gain
  - OSI 7 layer model, interfaces, protocols
  - Encapsulation, demultiplexing
- This lecture
  - Socket programming overview
  - Physical layer
- Online lecture videos: http://computernetworks5e.org

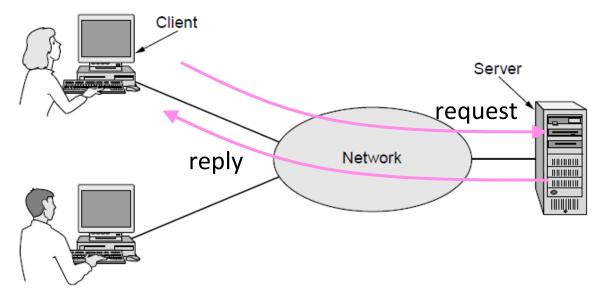
# **Network-Application Interface**

- Defines how apps use the network
  - Lets apps talk to each other via hosts; hides the details of the network



# **Motivating Application**

Simple client-server connection setup



# Motivating Application (2)

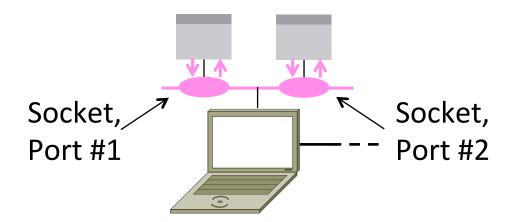
- Simple client-server connection setup
  - Client app sends a request to server app
  - Server app returns a (longer) reply
- This is the basis for many apps!
  - File transfer: send name, get file (§6.1.4)
  - Web browsing: send URL, get page
  - Echo: send message, get it back
- Let's see how to write this app ...

## Socket API

- Simple abstraction to use the network
  - The network service API used to write all Internet applications
  - Part of all major OSes and languages; originally Berkeley (Unix)
     ~1983
- Supports two kinds of network services
  - Streams: reliably send a stream of bytes
  - Datagrams: unreliably send separate messages. (Ignore for now.)

## Socket API (2)

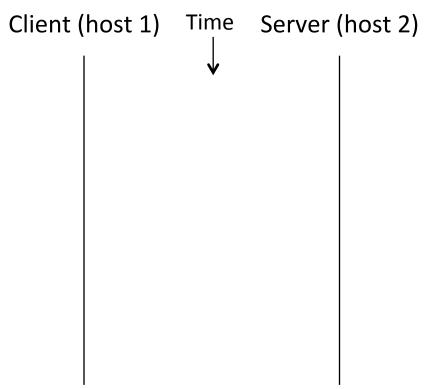
 <u>Sockets</u> let applications attach to the local network at different <u>ports</u>



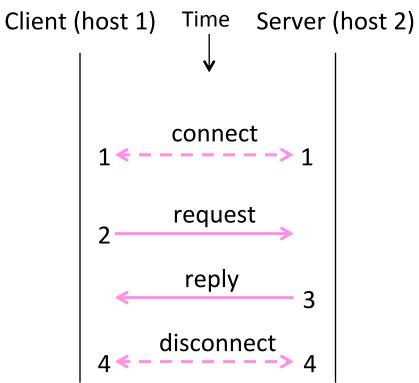
# Socket API (3)

Primitive	Meaning				
SOCKET	Create a new communication endpoint				
BIND	Associate a local address with a socket				
LISTEN	Announce willingness to accept connections; give queue size				
ACCEPT	Passively wait for an incoming connection				
CONNECT	Actively attempt to establish a connection				
SEND	Send some data over the connection				
RECEIVE	Receive some data from the connection				
CLOSE	Release the connection				

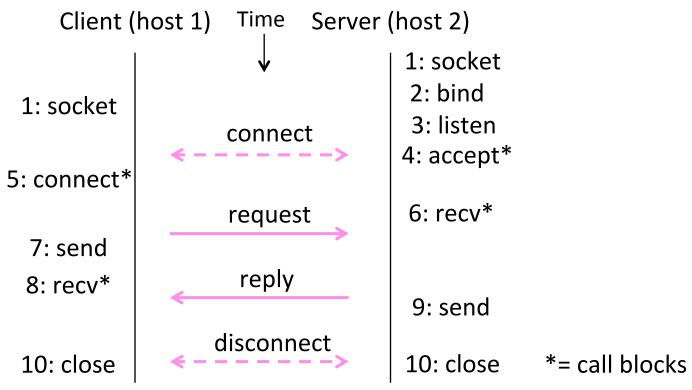
# **Using Sockets**



# Using Sockets (2)



# Using Sockets (3)



# Client Program (outline)

# Server Program (outline)

```
socket() // make socket
getaddrinfo() // for port on this host
bind() // associate port with socket
listen() // prepare to accept connections
accept() // wait for a connection [block]
...
recv() // wait for request
...
send() // send the reply
close() // eventually disconnect
```

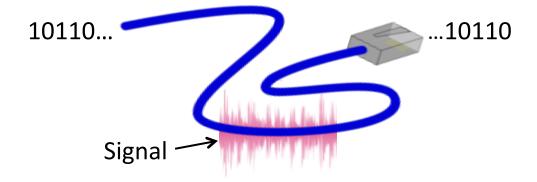
### Where we are in the Course

Beginning to work our way up starting with the Physical layer

Application
Transport
Network
Link
Physical

# Scope of the Physical Layer

- Concerns how signals are used to transfer message bits over a link
  - Wires etc. carry <u>analog signals</u>
  - We want to send <u>digital bits</u>

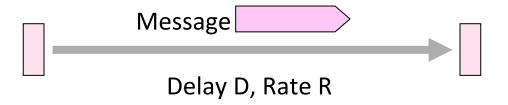


## **Topics**

- 1. Properties of media
  - Wires, fiber optics, wireless
- 2. Simple signal propagation
  - Bandwidth, attenuation, noise
- 3. Modulation schemes
  - Representing bits, noise
- 4. Fundamental limits
  - Nyquist, Shannon

# Simple Link Model

- We'll end with an abstraction of a physical channel
  - Rate (or bandwidth, capacity, speed) in bits/second
  - Delay or Latency in seconds, related to length



- Other important properties:
  - Whether the channel is broadcast, and its error rate

## Message Latency

- Latency L: delay to send a message over a link
  - Transmission delay: time to put M-bit message "on the wire"

```
T-delay = M (bits) / Rate (bits/sec) = M/R seconds
```

- Propagation delay: time for bits to propagate across the wire
  - P-delay = Length / speed of signals = Length / ⅓c = D seconds
- Combining the two terms we have: L = M/R + D

#### **Metric Units**

• The main prefixes we use:

Prefix	Exp.	prefix	exp.
K(ilo)	10 <sup>3</sup>	m(illi)	10 <sup>-3</sup>
M(ega)	10 <sup>6</sup>	μ(micro)	10 <sup>-6</sup>
G(iga)	10 <sup>9</sup>	n(ano)	10 <sup>-9</sup>

- Use powers of 10 for rates, 2 for storage or data size
  - 1 Mbps = 1,000,000 bps, 1 KB =  $2^{10}$  bytes
- "B" is for bytes, "b" is for bits

# Latency Examples

- "Dialup" with a telephone modem:
  - D = 5 ms, R = 56 kbps, M = 1250 bytes

- Broadband cross-country link:
  - D = 50 ms, R = 10 Mbps, M = 1250 bytes

# Latency Examples (2)

"Dialup" with a telephone modem:

D = 5 ms, R = 56 kbps, M = 1250 bytes  
L = 5 ms + 
$$(1250x8)/(56 \times 10^3)$$
 sec = 184 ms!

Broadband cross-country link:

```
D = 50 ms, R = 10 Mbps, M = 1250 bytes
L = 50 ms + (1250x8) / (10 x 10^6) sec = 51 ms
```

- A long link or a slow rate means high latency
  - Often, one delay component dominates

# **Bandwidth-Delay Product**

Messages take space on the wire!



 The amount of data in flight is the <u>bandwidth-delay (BD)</u> <u>product</u>

$$BD = R \times D$$

- Measure in bits, or in messages
- Small for LANs, big for "long fat" pipes

# Bandwidth-Delay Example

Fiber at home, cross-country

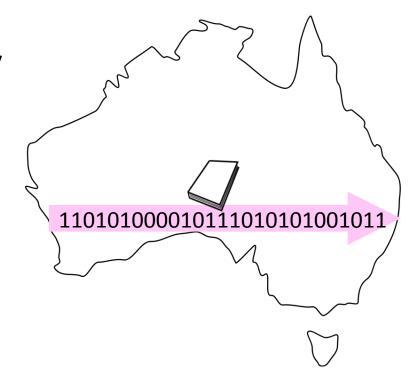
R=40 Mbps, D=50 ms

BD =  $40 \times 10^6 \times 50 \times 10^{-3}$  bits

= 2000 Kbit

= 250 KB

 That's quite a lot of data "in the network"!



# How "Long" is a Bit?

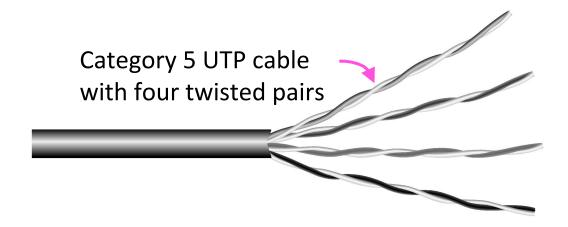
- Interesting trivia: how "long" is the representation of a bit on a wire?
- Considering a fiber optic cable
  - Signal propagation speed: 200'000'000 m/s
  - Sending rate: 1Gbps → duration of sending one bit: 1ns
  - Bit "length": 1ns \* 200'000'000 m/s = 0.2 m
  - "Length" of a 1Kb packet:  $0.2m * 8 * 2^{10} = 1.6km$

# Types of Media (§2.2, 2.3)

- Media propagate signals that carry bits of information
- We'll look at some common types:
  - Wires
  - Fiber (fiber optic cables)
  - Wireless

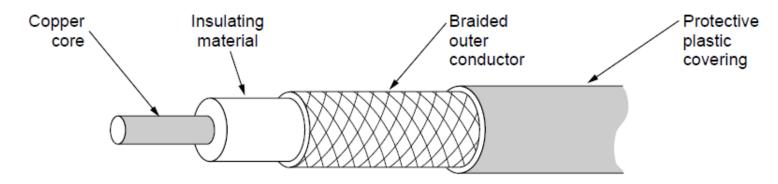
#### Wires – Twisted Pair

- Very common; used in LANs and telephone lines
  - Twists can reduce radiated signal or reduce effect of external interference signal



#### Wires – Coaxial Cable

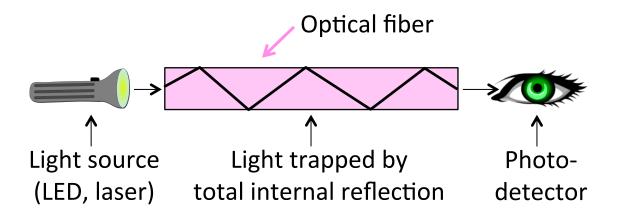
Also common. Better shielding for better performance



Other kinds of wires too: e.g., electrical power (§2.2.4)

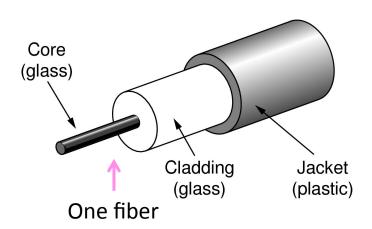
## Fiber

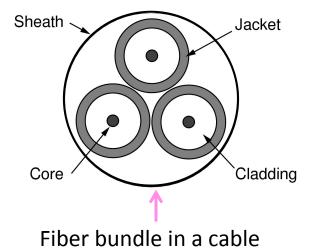
- Long, thin, pure strands of glass
  - Enormous bandwidth (high speed) over long distances



# Fiber (2)

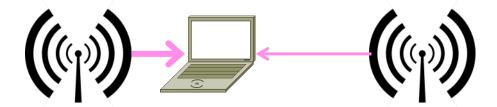
 Two varieties: multi-mode (shorter links, cheaper) and single-mode (up to ~100 km)





#### Wireless

- Sender radiates signal over a region
  - In many directions, unlike a wire, to potentially many receivers
  - Nearby signals (same freq.) <u>interfere</u> at a receiver; need to coordinate use



#### **UNITED**

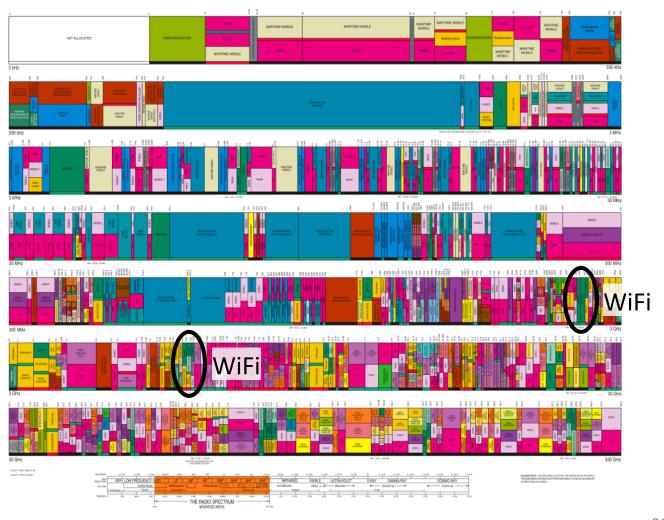
#### **STATES**

**FREQUENCY** 

#### **ALLOCATIONS**

#### THE RADIO SPECTRUM





NET	Ama Ama	FIXED SATELLITE  SATELLITE  SATELLITE  SATELLITE  SATELLITE  SATELLITE  SATELLITE  FIXED  FIXED	BROADCASTING **  HixeD  **  BROADCASTING  **  **  **  **  **  **  **  **  **	Sy MOBILE SATELLITE MOBILE FE-Sy Space Reservation Manual HADOL MANUE MANUEL HADOL MANUEL HADOL MANUEL MANUEL HADOL MANUEL	WUNGATION SATELLITE E-S) Satellite (E-S) Mobile SATELLITE FXED Lind Mobile SAT, (E-S) Mobile FX SAT, (E-S) Mobile Mobile SAT, (E-S) Mobile	Mobile SPACE RESEARCH (Fassive) AUTCAL RADIONAVIS, DN Space Res, (act.) CATION RADIONAVIS RADIONAVI
SATELLIFE (E-S)   SATELLIFE (F-S)   SATELLIFE	PADIOLOCATION 95.0  RADIO- LOCATION 100.0  RADIO- R	D6.0 ASTRONOMY	16.0 19.98 20.02 26.0	Kessaicii		151.0   WADILE
AMATEURS Amateur Amateur Sat. Amateur Sat. ApploLoc.  AMATEUR AMATEURS AMATEUR AMATEUR SATEUR	MOBILE CATION SATELLITE CATION SATELLITE SPACE RESEARCH PRESENCE FIXED F	EARTH EXPLORATION SATELLITE (Passive)  SPACE RESEARCH (Passive)	MOBILE SATELITE RESEARCH EVELSAT.  MO- INTER- SPACE EATH MOBILE SAT. RES. PP. 34  MOBILE SAT. RES. PP. 34  MOBILE SATELITE RESEARCH EVELSAT.  Passie)	INTER- SATELLITE  RADIO- LOCATION  RADIO- NAVIGATI  RADIO- NAVIGATI  MOBILE  MOBILE  MOBILE  MOBILE  MOBILE	AMATEUR SATELUTE	RADIO SPACE RES. STRONOMY (Passive) MOBILE SATELLITE

10.45 10.5 10.6

10.68

10.7

8.025 8.175 8.215

14.4 14.47 14.5 14.7145 15.1365

15.35

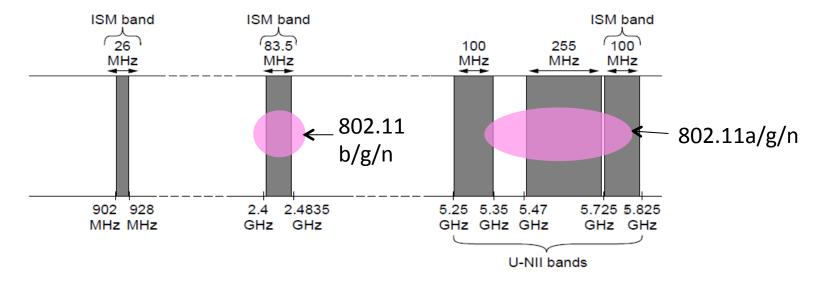
13.75

12.7

15.4 15.43 15.6 15.7 16.6

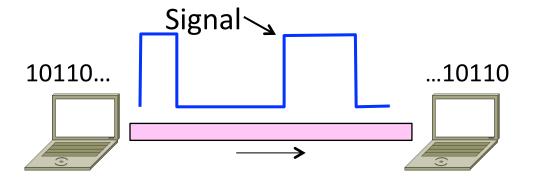
## Wireless (2)

 Microwave, e.g., 3G, and unlicensed (ISM: Industry Science Medicine) frequencies, e.g., WiFi, are widely used for computer networking



# Signals (§2.2, 2.3)

 Analog signals encode digital bits. We want to know what happens as signals <u>propagate</u> over media

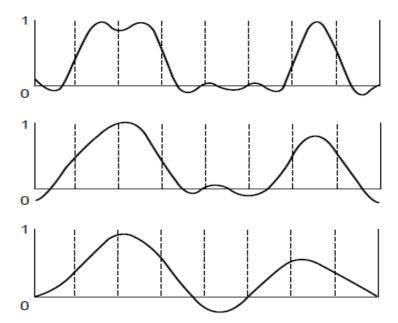


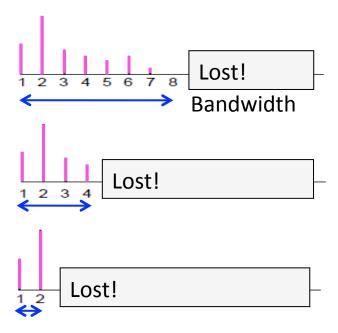
## Frequency Representation

 A signal over time can be represented by its frequency components (called Fourier analysis)

### Effect of Less Bandwidth

Fewer frequencies (=less bandwidth) degrades signal





### Signals over a Wire

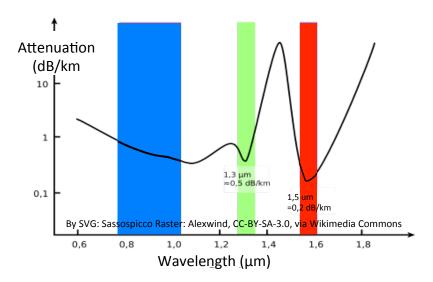
- What happens to a signal as it passes over a wire?
  - The signal is delayed (propagates at ¾c)
  - The signal is attenuated
  - 3. Frequencies above a cutoff are highly attenuated
  - 4. Noise is added to the signal (later, causes errors)

EE: Bandwidth = width of frequency band, measured in Hz

CS: Bandwidth = information carrying capacity, in bits/sec

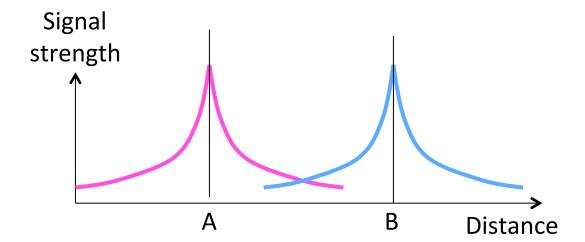
### Signals over Fiber

- Light propagates with very low loss in three very wide frequency bands
  - Use a carrier to send information



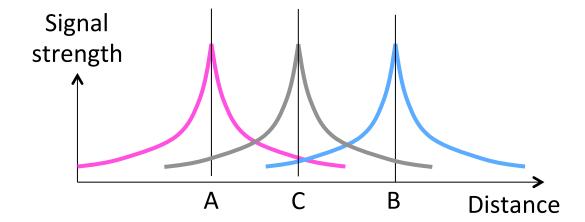
### Signals over Wireless

 Travel at speed of light, spread out and attenuate faster than 1/dist<sup>2</sup>



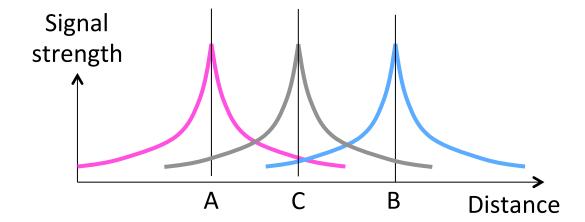
# Signals over Wireless (2)

Multiple signals on the same frequency interfere at a receiver



# Signals over Wireless (3)

Interference leads to notion of <u>spatial reuse</u> (of same freq.)

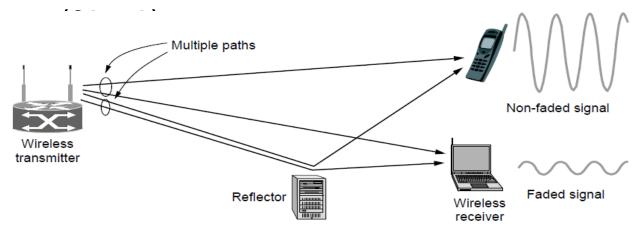


# Signals over Wireless (4)

- Various other effects too!
  - Wireless propagation is complex, depends on environment
- Some key effects are highly frequency dependent
  - E.g., <u>multipath</u> at microwave frequencies

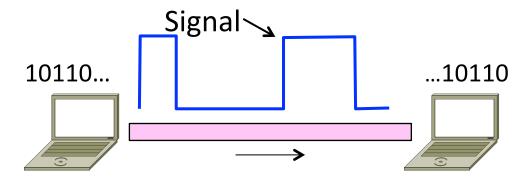
### Wireless Multipath

- Signals bounce off objects and take multiple paths
  - Some frequencies attenuated at receiver, varies with location
  - Messes up signal; handled with sophisticated methods



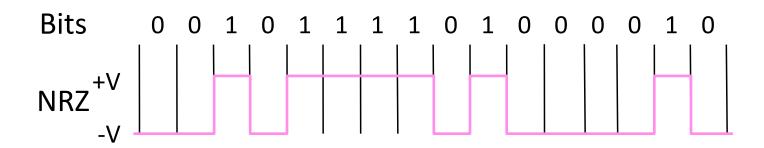
### Modulation (§2.5)

- We've talked about signals representing bits. How, exactly?
  - This is the topic of modulation



### A Simple Modulation

- Let a high voltage (+V) represent a 1, and low voltage (-V) represent a 0
  - This is called NRZ (Non-Return to Zero)



### Many Other Schemes

- Can use more signal levels, e.g., 4 levels is 2 bits per symbol
- Practical schemes are driven by engineering considerations
  - E.g., clock recovery

### **Clock Recovery**

- Um, how many zeros was that?
  - Receiver needs frequent signal transitions to decode bits

```
1 0 0 0 0 0 0 0 0 ... 0
```

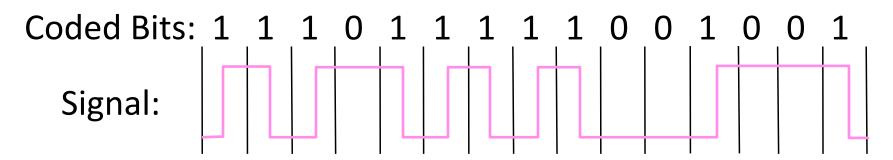
- Several possible designs
  - E.g., Manchester coding and scrambling (§2.5.1)

# Clock Recovery – 4B/5B

- Map every 4 data bits into 5 code bits without long runs of zeros
  - $-0000 \rightarrow 11110,0001 \rightarrow 01001,$  $1110 \rightarrow 11100, ... 1111 \rightarrow 11101$
  - Has at most 3 zeros in a row
  - Also invert signal level on a 1 to break up long runs of 1s (called NRZI, §2.5.1)

# Clock Recovery – 4B/5B (2)

- 4B/5B code for reference:
  - $-0000 \rightarrow 11110,0001 \rightarrow 01001,1110 \rightarrow 11100,...1111 \rightarrow 11101$
- Message bits: 1111 0000 0001



#### **Passband Modulation**

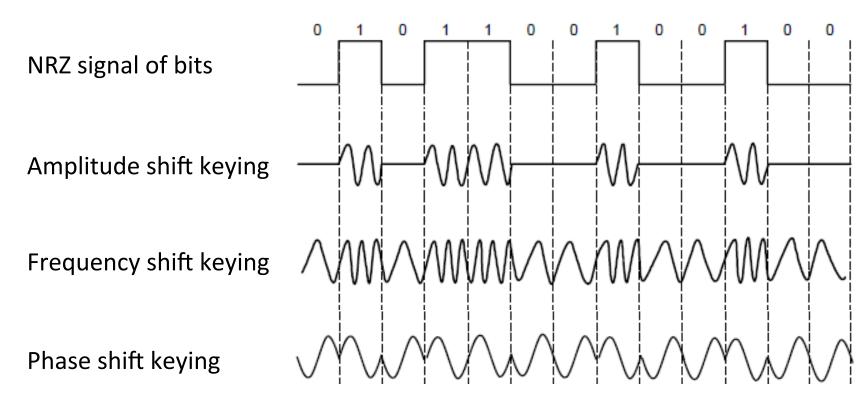
- What we have seen so far is <u>baseband</u> modulation for wires
  - Signal is sent directly on a wire
- These signals do not propagate well on fiber / wireless
  - Need to send at higher frequencies
- <u>Passband</u> modulation carries a signal by modulating a carrier

### Passband Modulation (2)

 Carrier is simply a signal oscillating at a desired frequency:

- We can modulate it by changing:
  - Amplitude, frequency, or phase

### Passband Modulation (3)



### Fundamental Limits (§2.1)

- How rapidly can we send information over a link?
  - Nyquist limit (~1924)
  - Shannon capacity (1948)
- Practical systems are devised to approach these limits

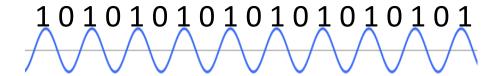
### **Key Channel Properties**

- The bandwidth (B), signal strength (S), and noise strength (N)
  - B limits the rate of transitions
  - S and N limit how many signal levels we can distinguish



### **Nyquist Limit**

The maximum <u>symbol</u> rate is 2B



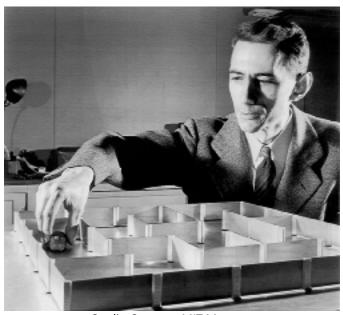
 Thus if there are V signal levels, ignoring noise, the maximum bit rate is:

$$R = 2B \log_2 V bits/sec$$

# Claude Shannon (1916-2001)

- Father of information theory
  - "A Mathematical Theory of Communication", 1948
- Fundamental contributions to digital computers, security, and communications

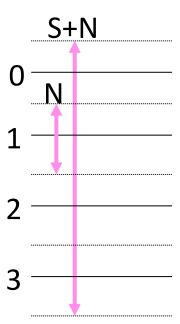
Electromechanical mouse that "solves" mazes!



Credit: Courtesy MIT Museum

# **Shannon Capacity**

- How many levels we can distinguish depends on S/N
  - Or SNR, the <u>Signal-to-Noise Ratio</u>
  - Note noise is random, hence some errors
- SNR given on a log-scale in deciBels:
  - $-SNR_{dB} = 10log_{10}(S/N)$



# **Shannon Capacity (2)**

• Shannon limit is for capacity (C), the maximum information carrying rate of the channel:

 $C = B \log_2(1 + S/N)$  bits/sec

# Wired/Wireless Perspective

- Wires and Fiber
- Engineer SNR for data rate
- Engineer link to have requisite SNR and B
- →Can fix data rate
- Wireless

Adapt data rate to SNR

- Given B, but SNR varies greatly, e.g., up to 60 dB!
- →Can't design for worst case, must adapt data rate

# Putting it all together – DSL

- DSL (Digital Subscriber Line, see §2.6.3) is widely used for broadband; many variants offer 10s of Mbps
  - Reuses twisted pair telephone line to the home; it has up to
     MHz of bandwidth but uses only the lowest ~4 kHz







### **DSL** (2)

- DSL uses passband modulation (called OFDM §2.5.1)
  - Separate bands for upstream and downstream (larger)
  - Modulation varies both amplitude and phase (called QAM)
  - High SNR, up to 15 bits/symbol, low SNR only 1 bit/symbol

