Design of Parallel and High-Performance Computing

Fall 2014

Lecture: Scheduling

Instructor: Torsten Hoefler & Markus Püschel

TA: Timo Schneider & Arnamoy Bhattacharyya

ETH

Eidgenössische Technische Hochschule Zürich
Swiss Federal Institute of Technology Zurich

Overview

- **DAGs again: An example**
- **Scheduling**
  - Greedy
  - Work stealing
- **Cilk**

- **Background material:**
  - Blumofe, Leiserson, Scheduling Multithreaded Computations by Work Stealing, Journal ACM, 46(5), 1999

Example: Fibonacci Numbers

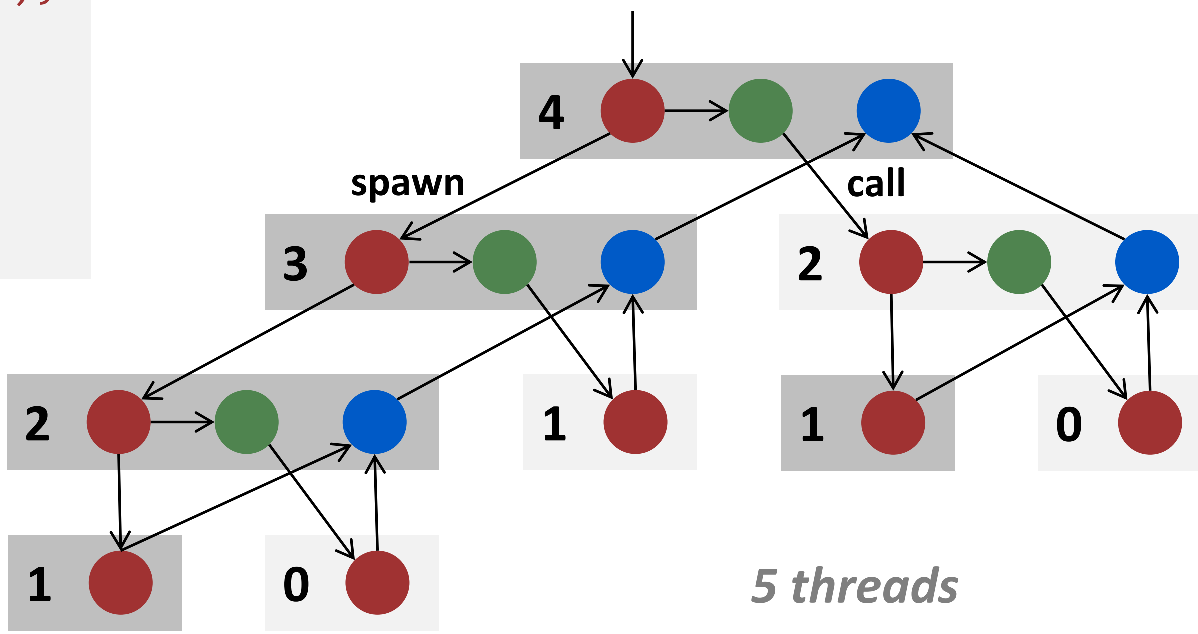
```
int fib (int n) {
  if (n<2) return (n);
  else {
    int x,y;
    x = spawn fib(n-1); // can execute in
                        // parallel with parent
    y = fib(n-2);
    sync;
    return (x+y);
  }
}
```

Stupid way of computing (why?)
But good example

Example: Fibonacci Numbers

```
int fib (int n) {
  if (n<2) return (n);
  else {
    int x,y;
    x = spawn fib(n-1);
    y = fib(n-2);
    sync;
    return (x+y);
  }
}
```

The DAG unfolds dynamically:



5 threads

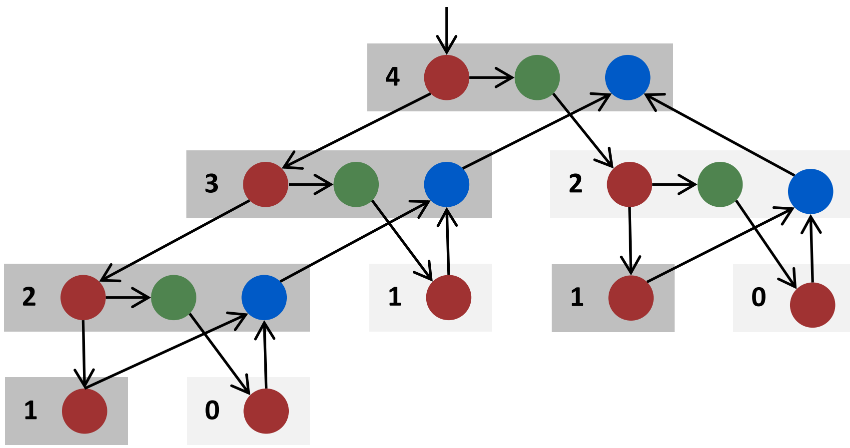
Node: **Sequence of instructions without call, spawn, sync, return**

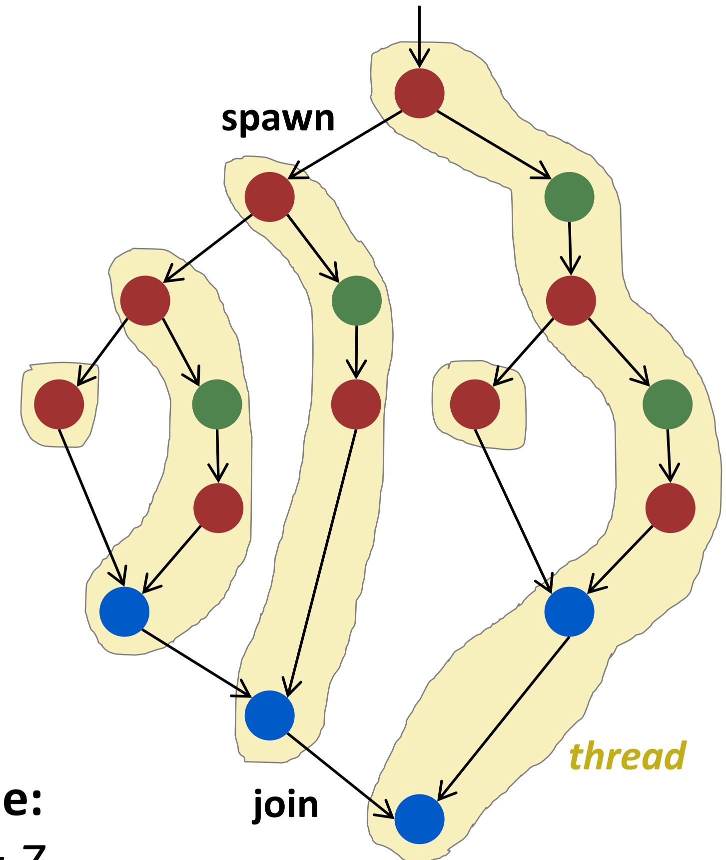
Edge: **Dependency**

Example: Fibonacci Numbers

Graphs obtained this way are called nested parallel (or fully strict):

- Every thread has one incoming edge (the spawn edge)
- All join edges from a thread connected to the parent thread

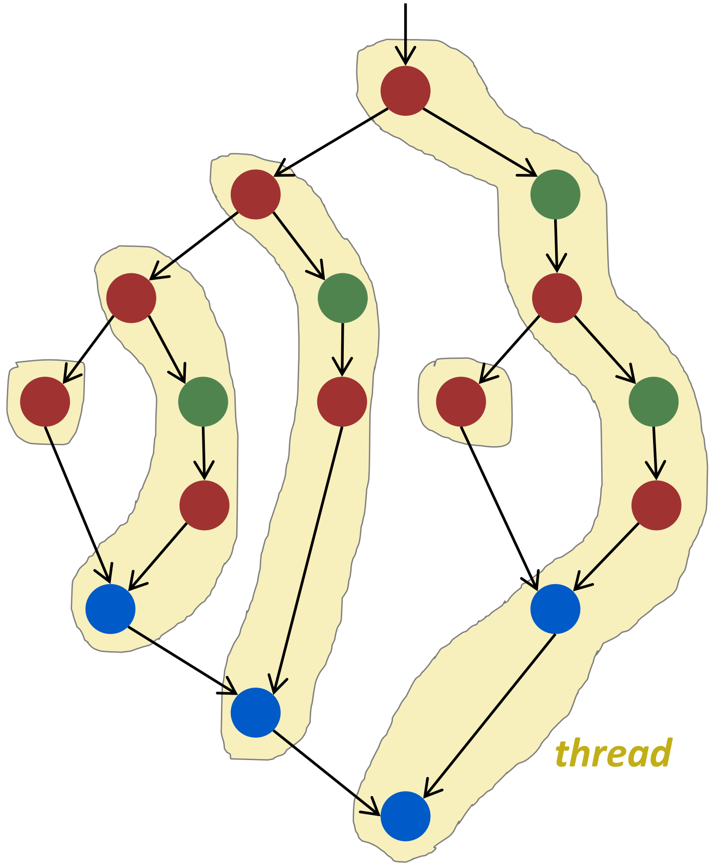




Assuming every node has unit time:

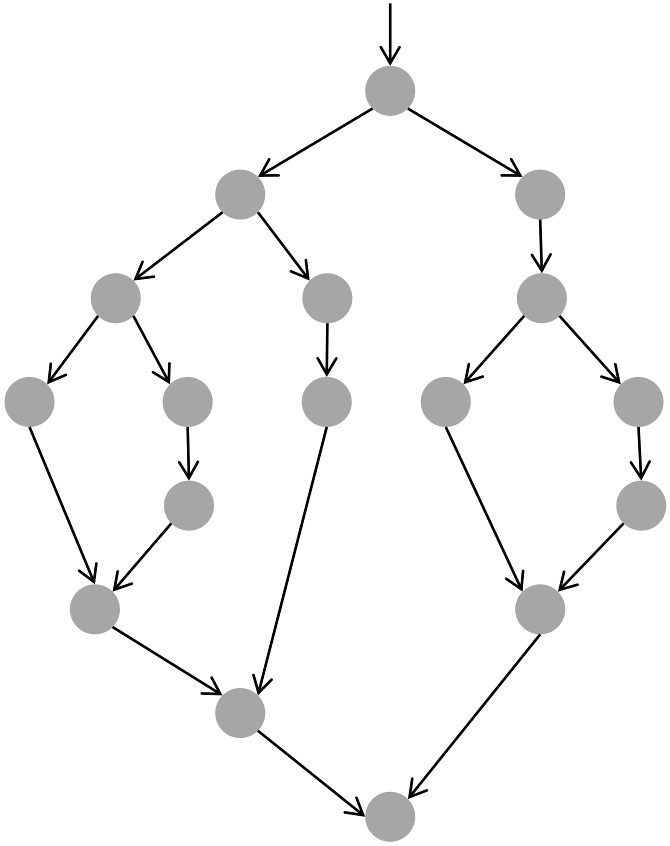
W = 17, D = 7

How to Schedule on p Processors?



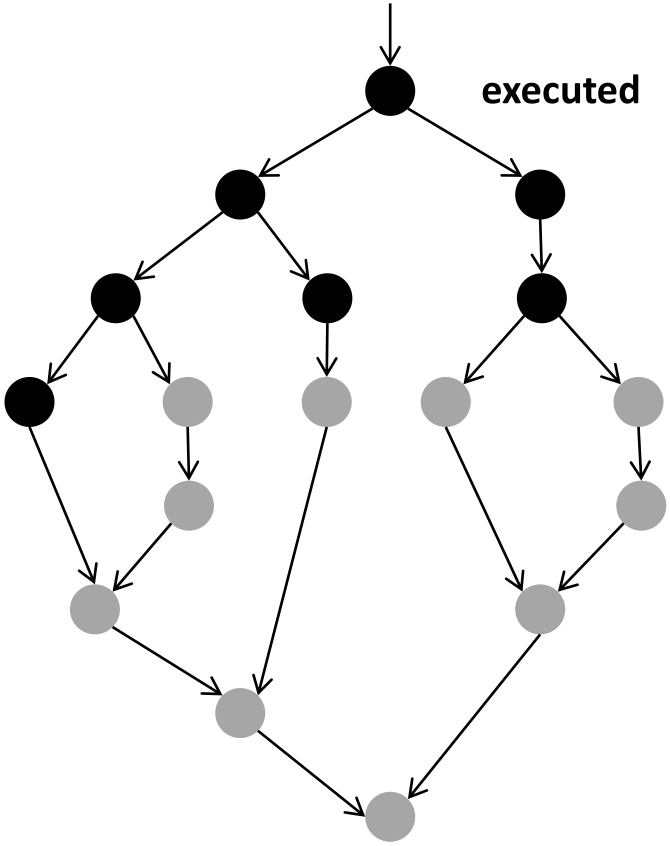
Greedy Scheduler

- *Idea:* Do as much as possible in every step



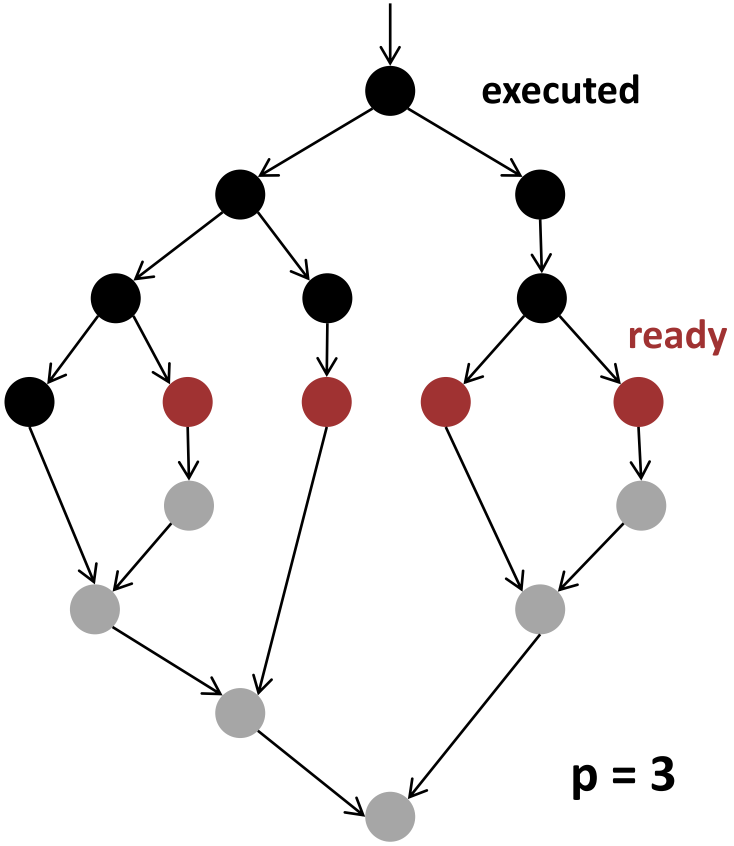
Greedy Scheduler

- *Idea:* Do as much as possible in every step
- *Definition:* A node is ready if all predecessors have been executed



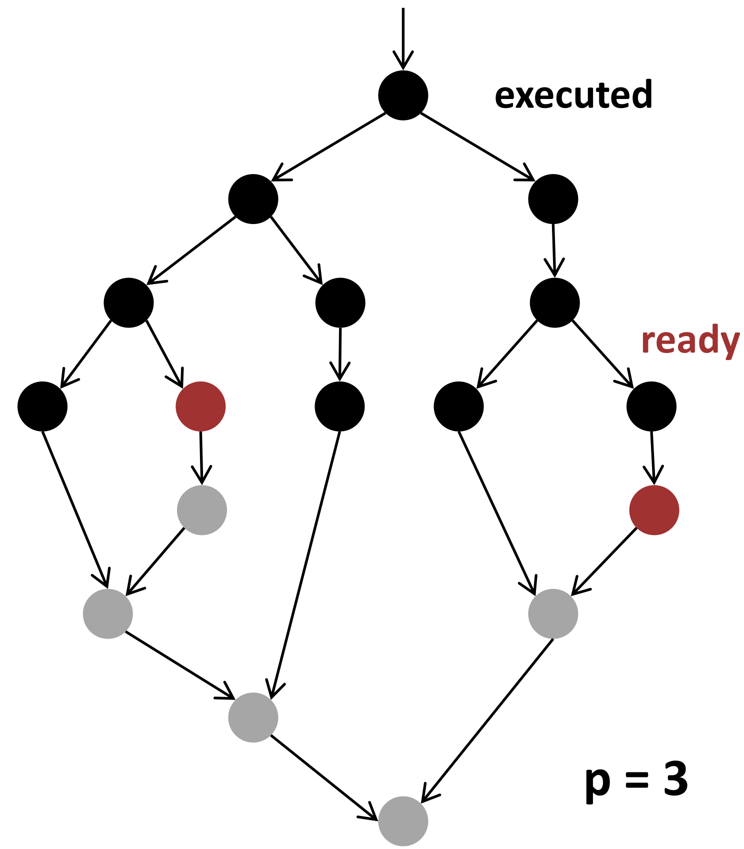
Greedy Scheduler

- *Idea:* **Do as much as possible in every step**
- *Definition:* **A node is ready if all predecessors have been executed**
- **Complete step:**
  - ≥ p nodes are ready
  - run any p



Greedy Scheduler

- *Idea:* **Do as much as possible in every step**
- *Definition:* **A node is ready if all predecessors have been executed**
- **Complete step:**
  - ≥ p nodes are ready
  - run any p
- **Incomplete step:**
  - < p nodes ready
  - run all
- **How good is this theoretically?** *(blackboard)*



Greedy Scheduler: Sketch

Maintain thread pool of live threads, each is ready or not

- **Initial: Root thread in thread pool, all processors idle**
- **At the beginning of each step each processor is idle or has a thread T to work on**
- **If idle**
  - *Get ready thread from pool*
- **If has thread T**
  - Case 0: T has another instruction to execute
    execute it
  - Case 1: thread T spawns thread S
    return T to pool, continue with S
  - Case 2: T stalls
    return T to pool, then idle
  - Case 3: T dies
    if parent of T has no living children, continue with the parent, otherwise idle

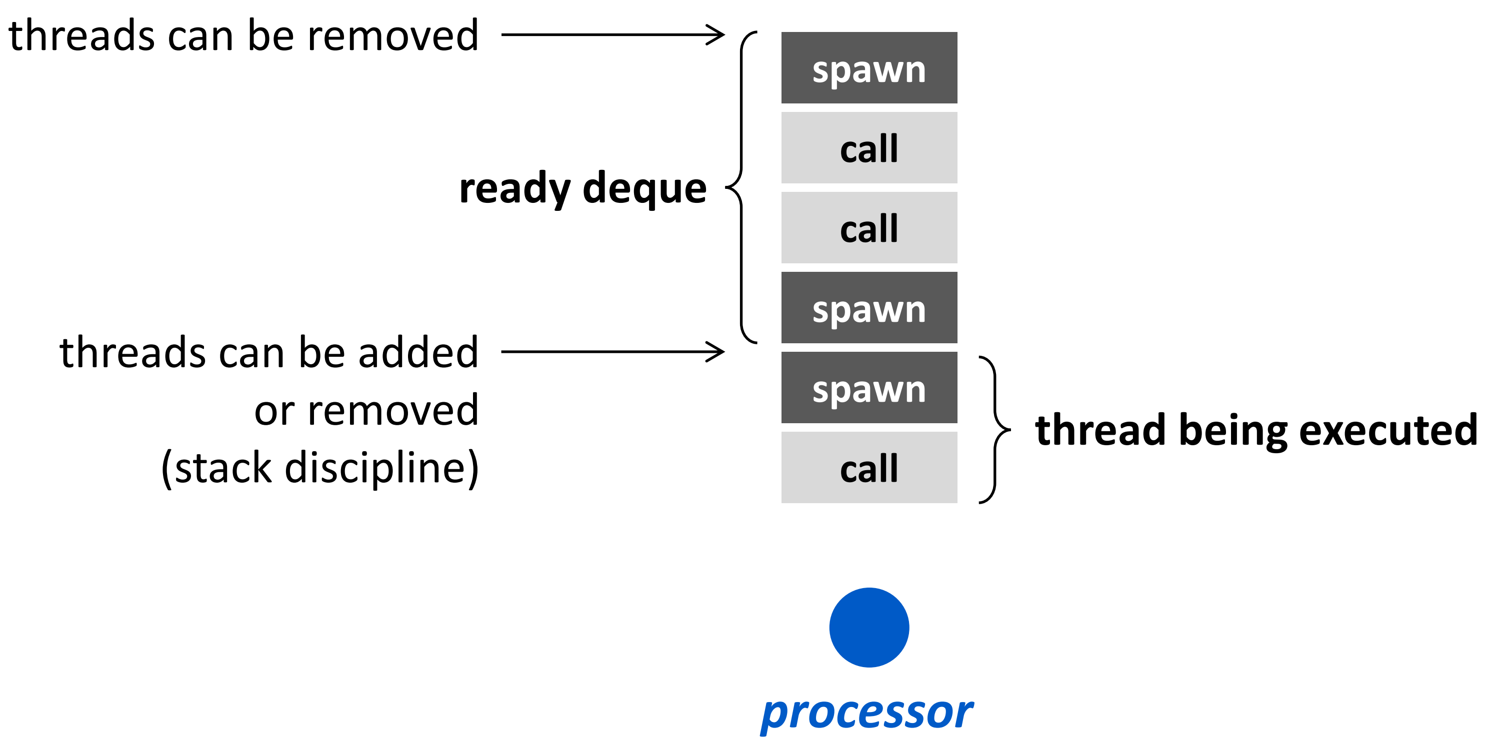
Greedy Scheduler: Problems

- **Centralized**
- **Overhead**

- **Work stealing scheduler:**
  - thread pool distributed
  - all processors do only useful work and operate locally as long as there is work to do
  - Good asymptotic behavior, good practical behavior
  - Implemented in Cilk runtime system

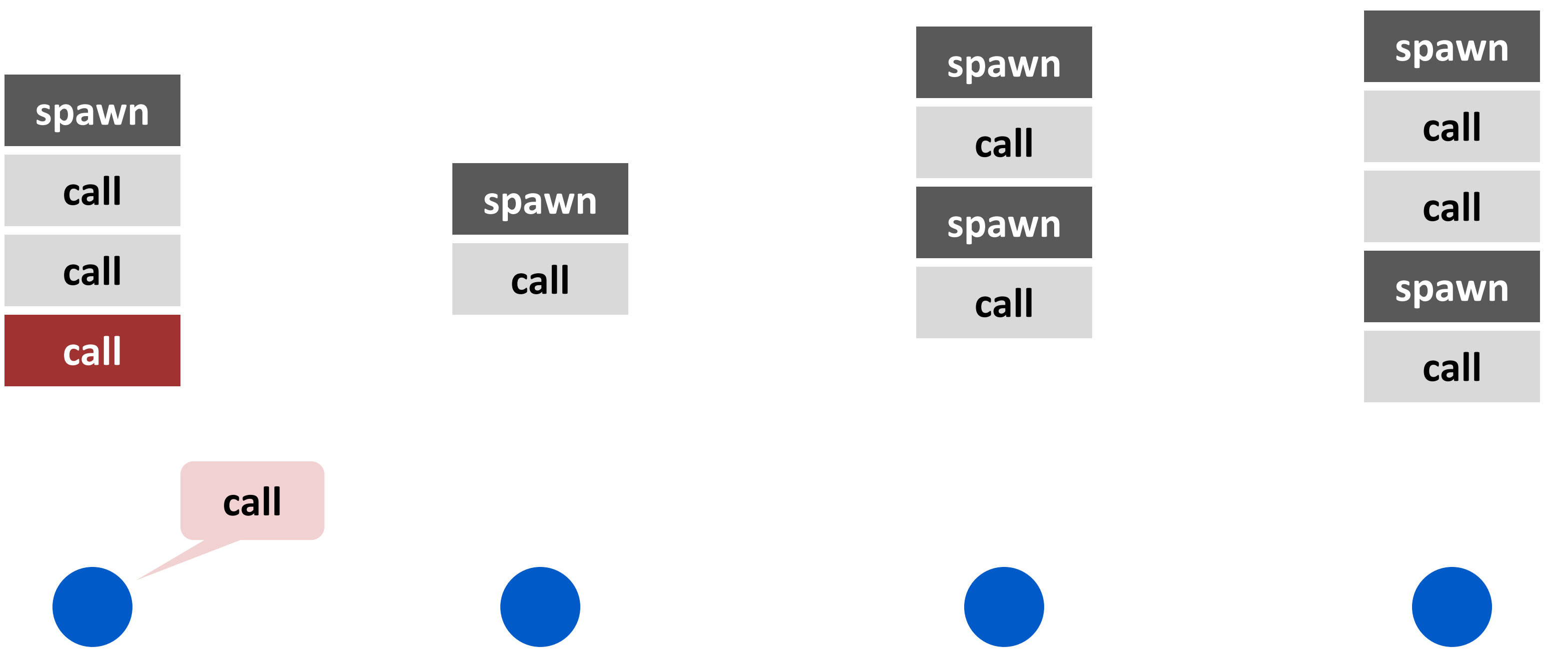
Work Stealing Scheduler

- **Each processor maintains a "ready deque:" deque of threads ready for execution; bottom is manipulated as a stack**



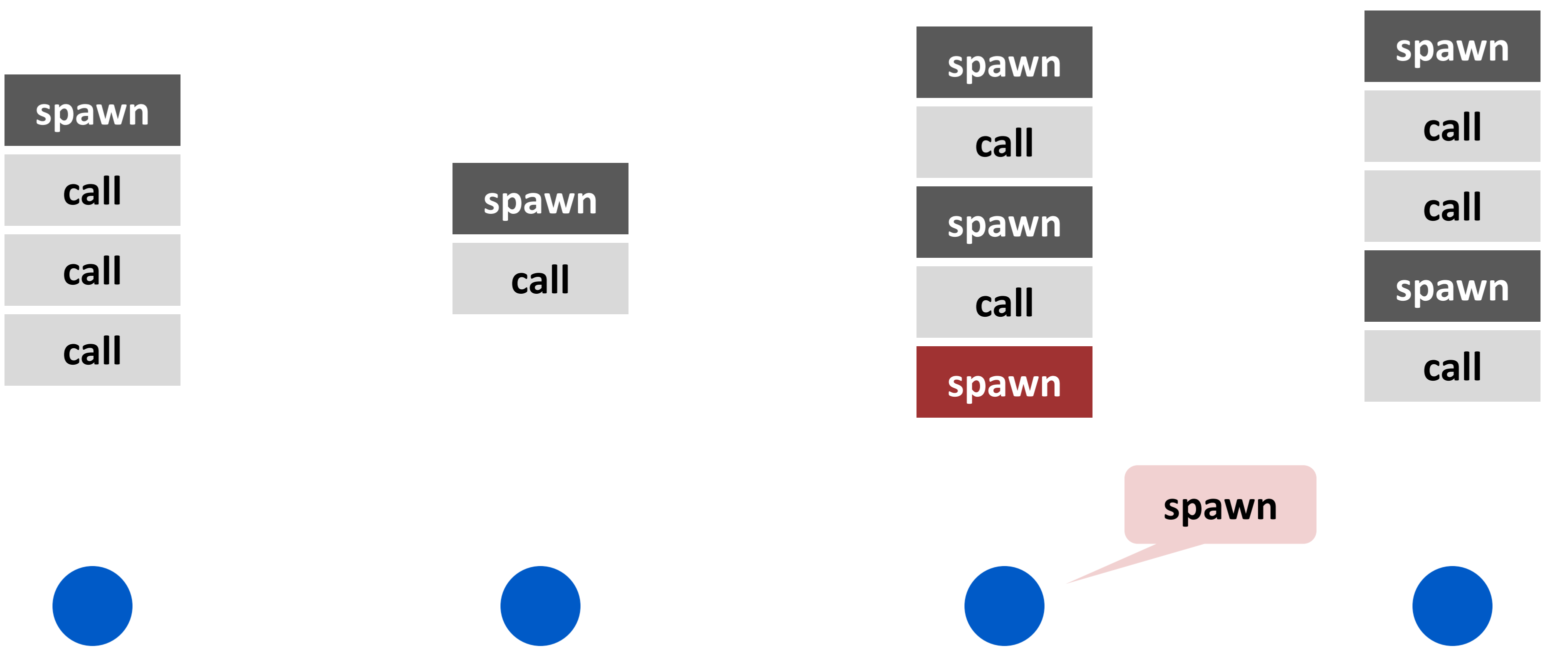
Work Stealing Scheduler

- Each processor maintains a “ready deque:” deque of threads ready for execution; bottom is manipulated as a stack



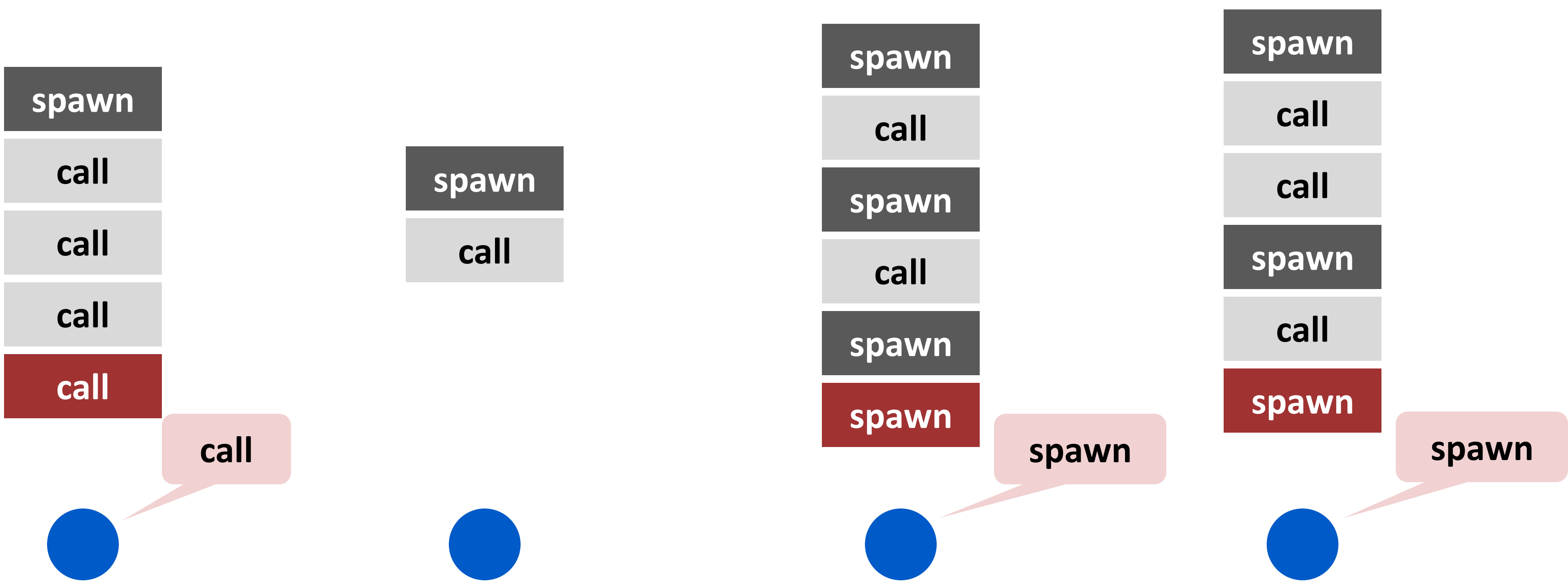
Work Stealing Scheduler

- Each processor maintains a “ready deque:” deque of threads ready for execution; bottom is manipulated as a stack



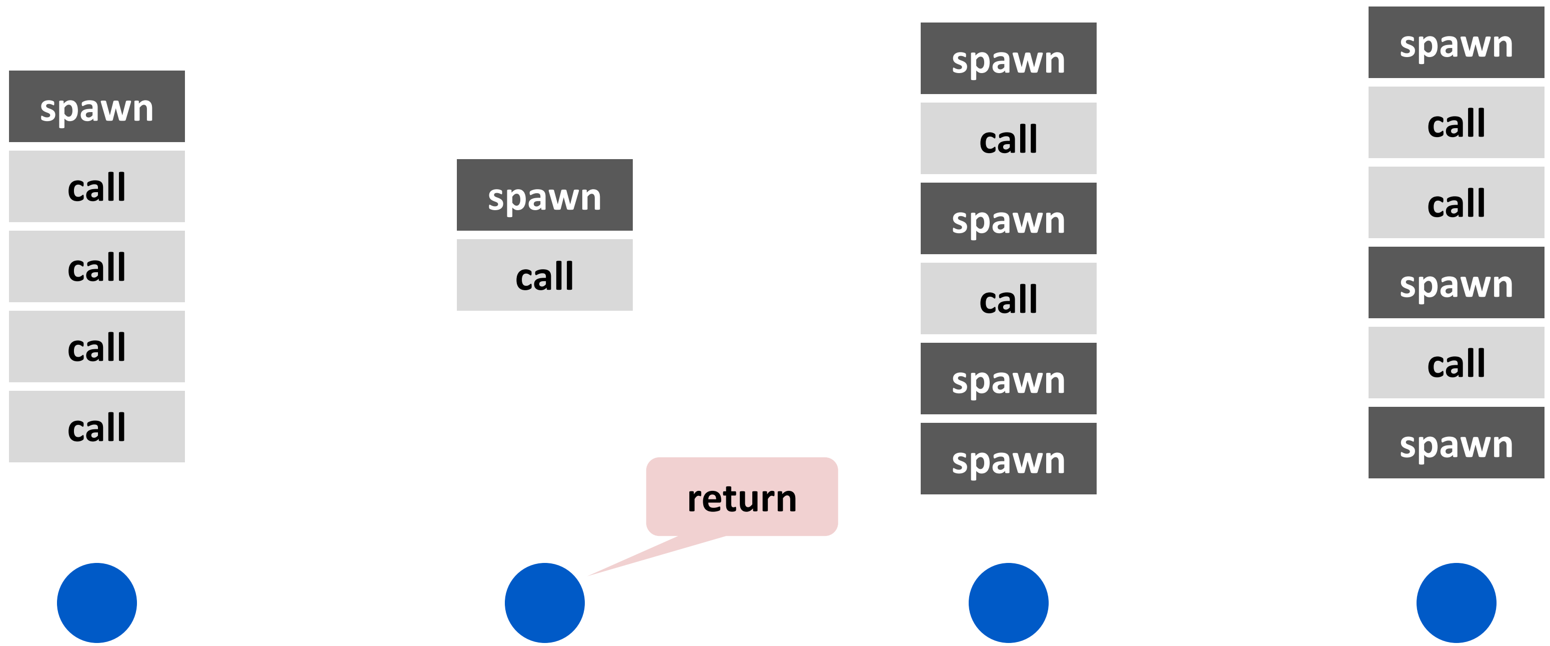
Work Stealing Scheduler

- **Each processor maintains a "ready deque:" deque of threads ready for execution; bottom is manipulated as a stack**



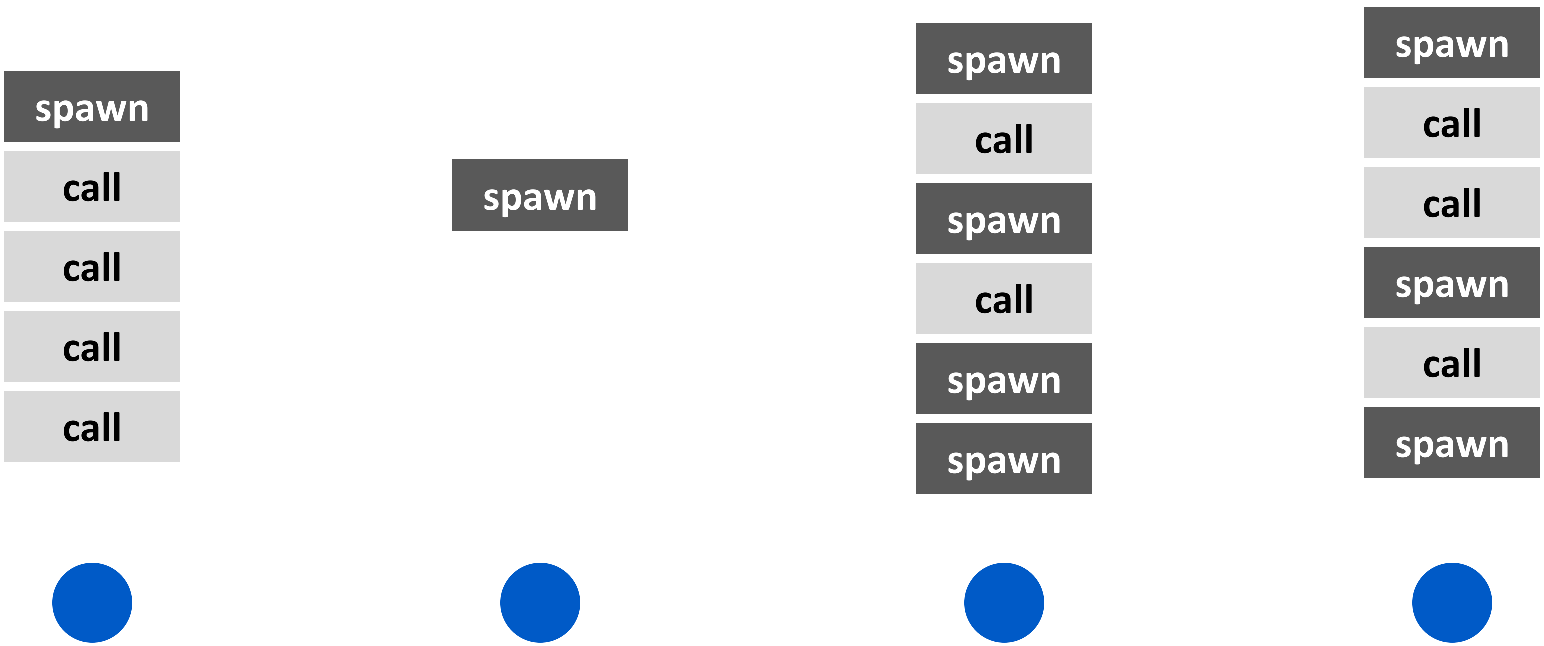
Work Stealing Scheduler

- Each processor maintains a “ready deque:” deque of threads ready for execution; bottom is manipulated as a stack



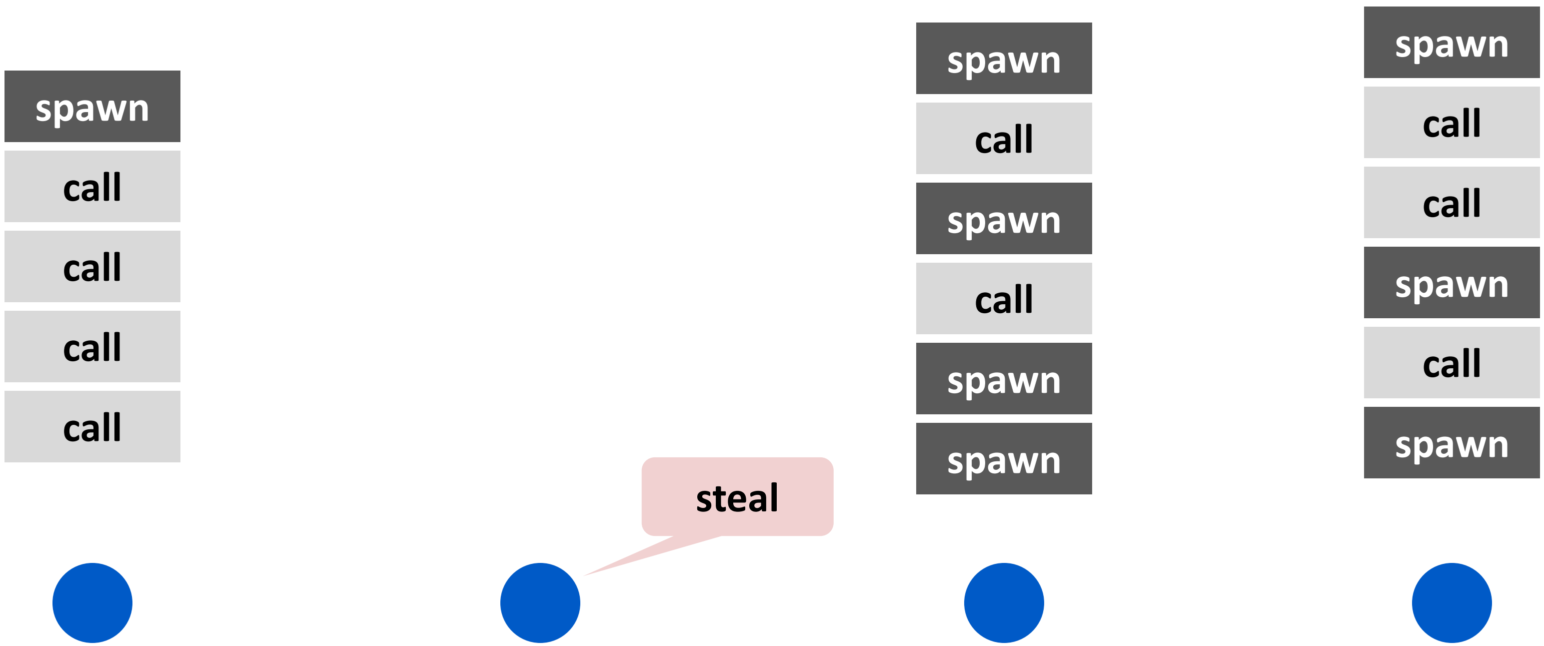
Work Stealing Scheduler

- Each processor maintains a “ready deque:” deque of threads ready for execution; bottom is manipulated as a stack



Work Stealing Scheduler

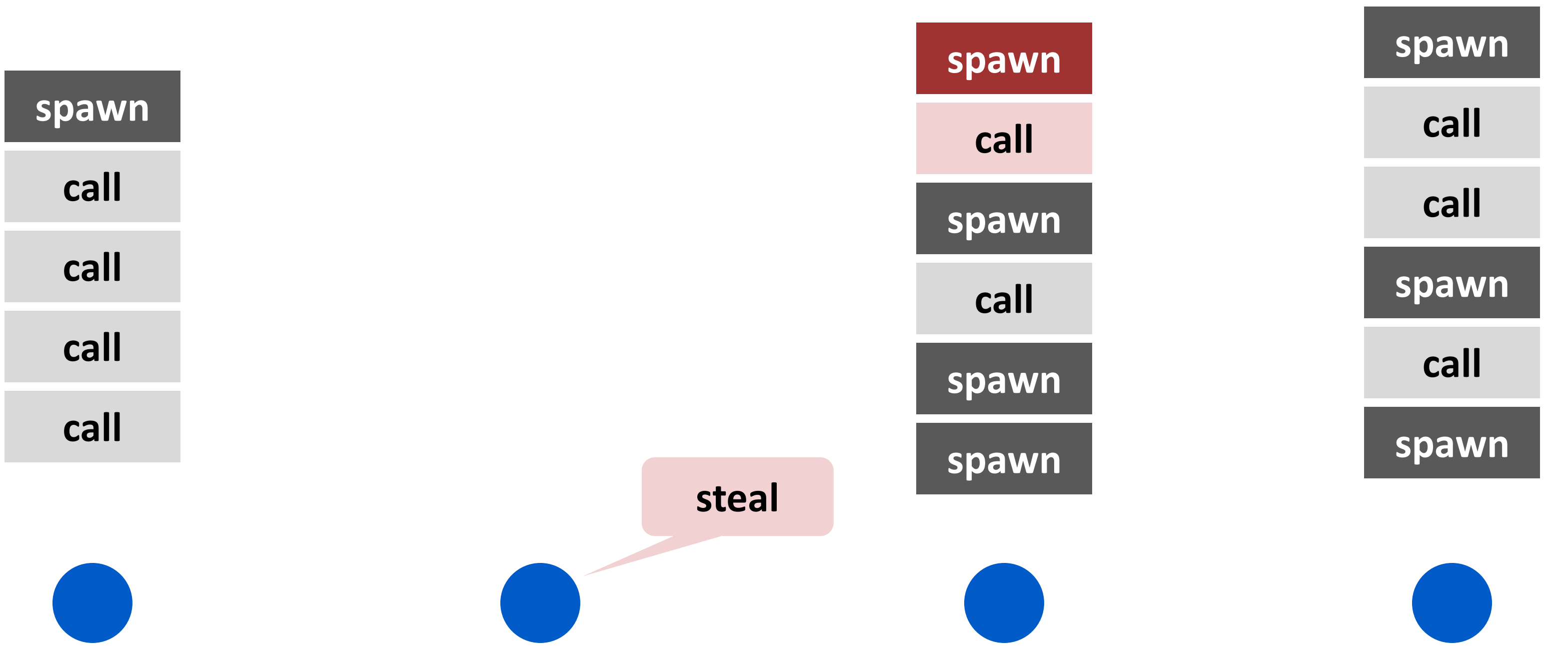
- Each processor maintains a "ready deque:" deque of threads ready for execution; bottom is manipulated as a stack



- Steal from the top of a randomly selected processor

Work Stealing Scheduler

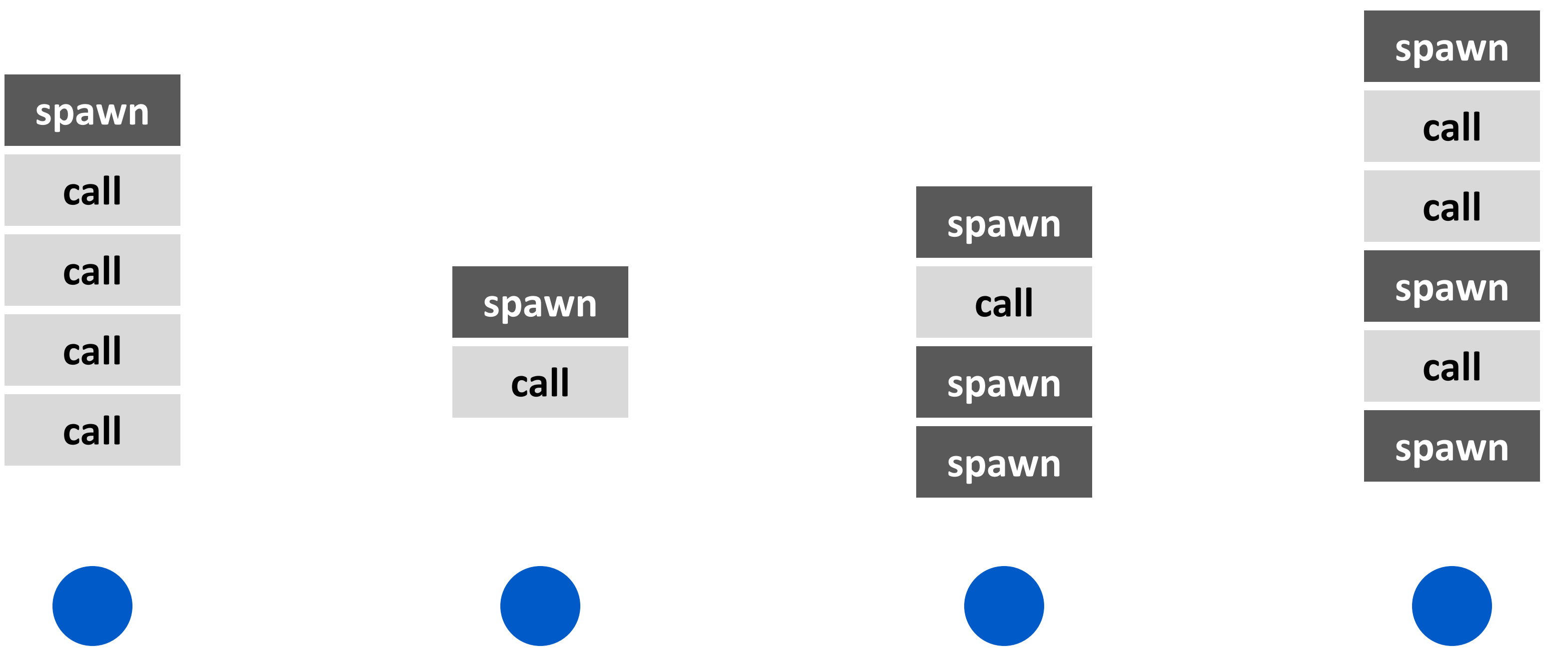
- Each processor maintains a “ready deque:” deque of threads ready for execution; bottom is manipulated as a stack



- Steal from the top of a randomly selected processor

Work Stealing Scheduler

- Each processor maintains a "ready deque:" deque of threads ready for execution; bottom is manipulated as a stack



Work Stealing Scheduler: Sketch

Each processor maintains a ready deque, bottom treated as stack

- **Initial: Root thread in deque of a random processor**
- **Deque not empty:**
  - Processor takes thread T from bottom and starts working
  - T spawns S: Put T on stack, continue with S
  - T stalls: Take next thread from stack
  - T dies: Take next thread from stack
  - If T enables a stalled thread S, S is put on the stack of T's processor
- **Deque empty:**
  - Steal thread from the top of a random (uniformly) processor's deque

- **Theoretical performance?** *(blackboard)*

Cilk

- Extension of C/C++
- Compiler and runtime system
- Developed at MIT, now distributed by Intel
- Cilk home at Intel