Design of Parallel and High Performance Computing Fall 2013

About projects

Instructors: Torsten Hoefler & Markus Püschel TA: Timo Schneider

ETH

Eidgenössische Technische Hochschule Zürich Swiss Federal Institute of Technology Zurich

Project: Rules

- Count 50% of the grade (work, presentation, report)
- Teams of two
 - Important: organize yourselves
- Topic: Some suggestions in a minute
- Timeline:
 - End Oct: Present your project in recitations
 - Late Nov/early Dec: Possibly progress presentations
 - Last week of class: Final project presentations
- Report:
- 6 pages
 - template provided
 - due January

Projects: Performance Optimization

- Pick an important algorithm/application
- Develop a parallel implementation that scales well on multicore
- Includes thorough benchmarking and experimental evaluation
- Requirements:
 - No numerical algorithm (dominated by floating point operations) Exceptions possible if directly related to student's research
 - Not sorting or anything that is mainly sorting

Example From Before



Example From Before

 Uses our fastest implementations depending on input size and adapts #threads accordingly

Bitonic Mergesort SSE	L Rad	SD ixsort	Parallel Radixsort with SSE		Pa Ra	arallel dixsort	Parallel Bitonic Mergesort SSE		
-									
2 ⁸	2 ¹⁰	2 ¹²	2 ¹⁴	2 ¹⁶	2 ¹⁸	2 ²⁰	2 ²²	2 ²⁴	2 ²⁶
				Input S	ize				
				mput	nic -				

Project Proposals

Advisor: Torsten Hoefler TA: Timo Schneider

Parallel Priority Queue (I)

Maintain a collection of data items, identified by a key. Finding the k smallest items (with the k smallest keys) should be supported on O(k) time. Finding any item by key should also be supported.

Required Operations

- queue_t init()
- void insert(queue_t q, void* data, uint64_t key)
- void*find(queue_t q, uint64_t key)
- void delete(queue_t q, uint64_t key)
- void*pop_front(queue_t q) // returns smallest element
- void finalize(queue_t q)

Collective Communications

- Assume P threads in shared memory
- Each thread p has:
 - a set of input elements i_{j,p} (0≤j<n-1)
 - a set of output elements o_{j,p} (0≤j<n-1)
- The post-condition (result) is:
 - $o_{j,p} = \sum_{p=1}^{P} i_{j,p} (0 \leq j < n)$. i.e., all $o_{\rm j,p}$ are identical on all p
- Tips:
 - Use the memory hierarchy and CC protocols (inline assembly is allowed!)
 - First optimize small n, then large n

Parallel Priority Queue (II)

Requirements contd.

- Multiple threads will be accessing the queue simultaneously (with all operations)
- Code may be written in C/C++ (gcc inline assembly is allowed ;-))
- Tips:
 - Experiment with different locking strategies and compare the performance
 - Pay attention to larger number of threads

Parallel BFS

- Generate an ER graph G(n,p) given n and p
- Perform a breath first search from n/2 vertices
 - Print the average maximum distance for any vertex
- Your implementation should exploit all available cores and perform the BFS as fast as possible

Parallel Graph Algorithms

Many more!

- Connected Components (CC)
- SSSP
- APSP (maybe too simple, looks like MatVec)
- Minimum spanning tree (MST)
- Vertex coloring
- Strongly connected components
- ... pick one and enjoy!
- Others
 - A* search
 - Various ML and AI algorithms (only nontrivial ones)

Schedule

- Some recitations will be used to demonstrate concepts in practice
 E.g., OpenMP basics, MPI basics, ...
- We will discuss "how to measure and report performance"
 - This is a complex topic often done wrong